



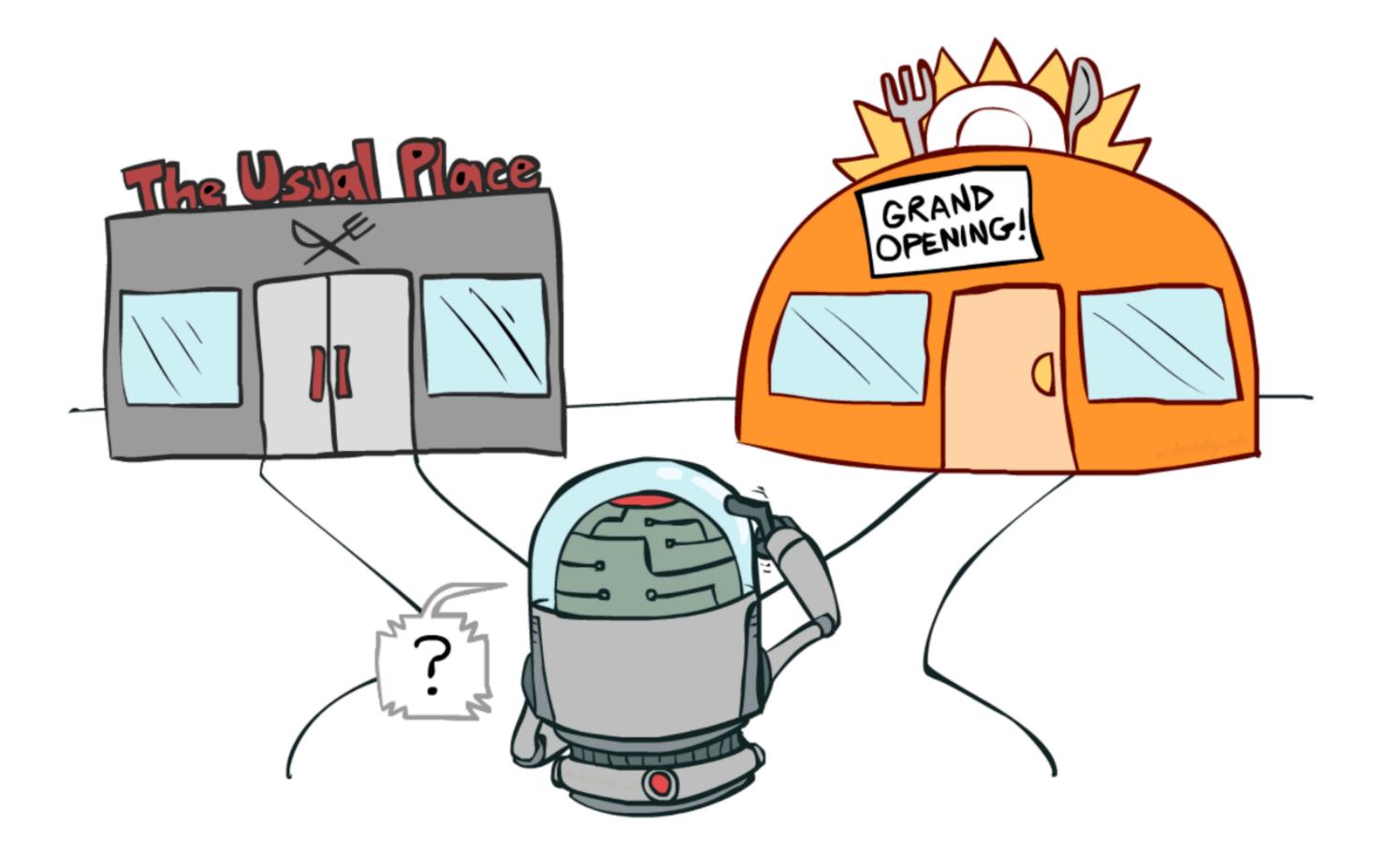
# **Exploration-Exploitation Dilemma** in Multi-Armed Bandit

Kimang KHUN, Ph.D.

Ministry of Industry, Science, Technology & Innovation

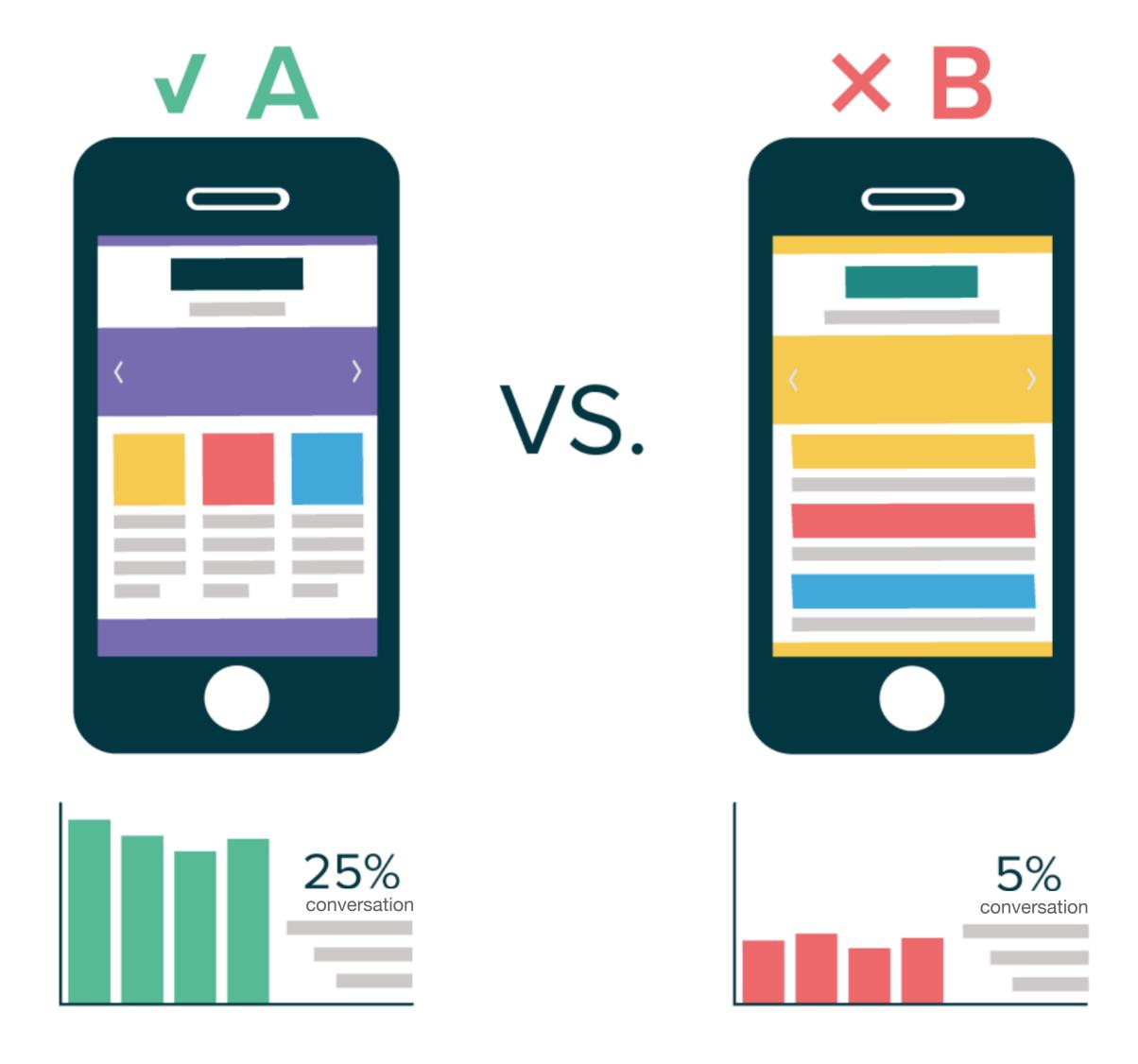
Seminar of AMS Department, Institute of Technology of Cambodia, Phnom Penh

## Exploration and Exploitation Dilemma

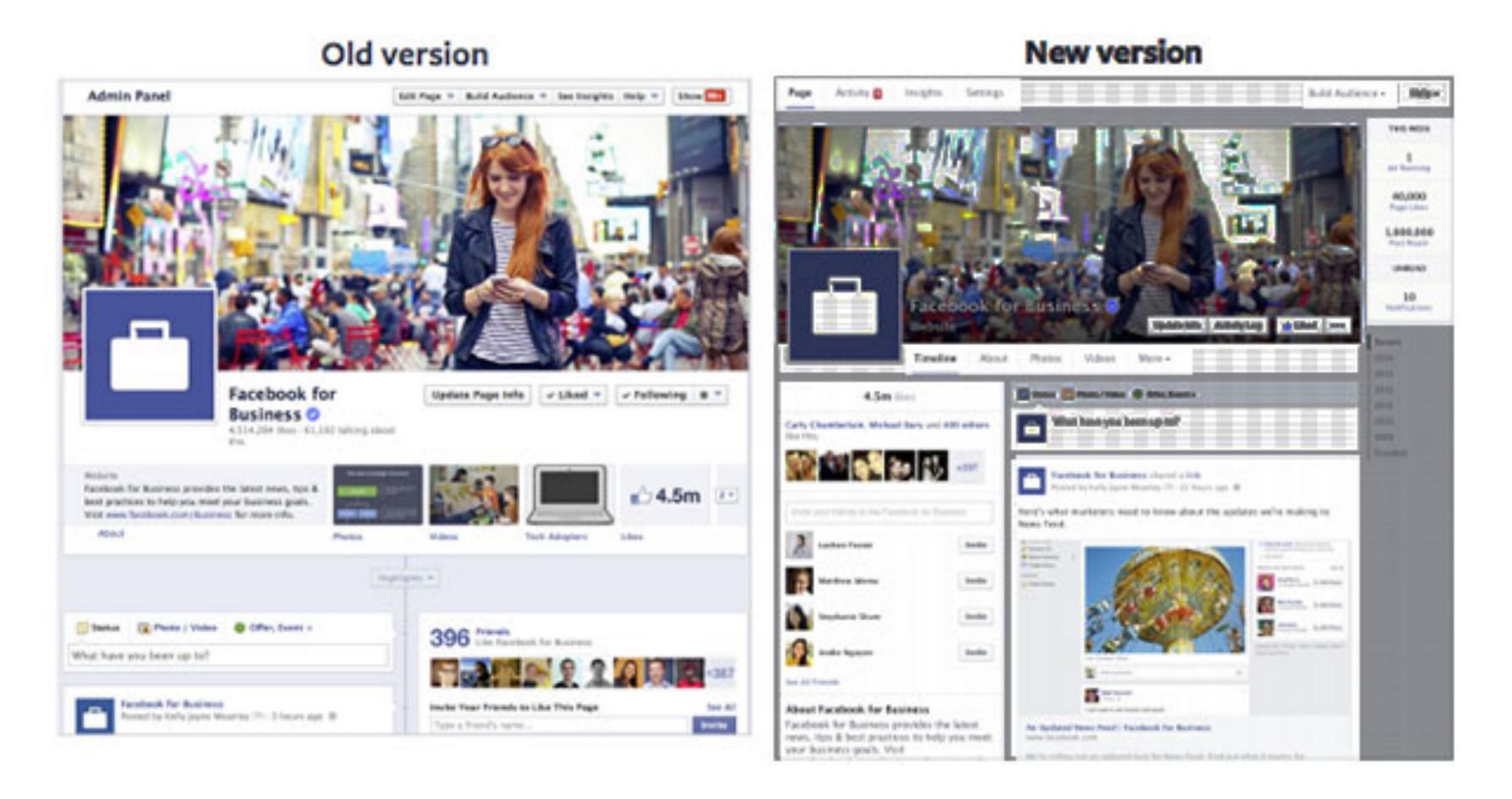


Source: Nicolas Gast's slides

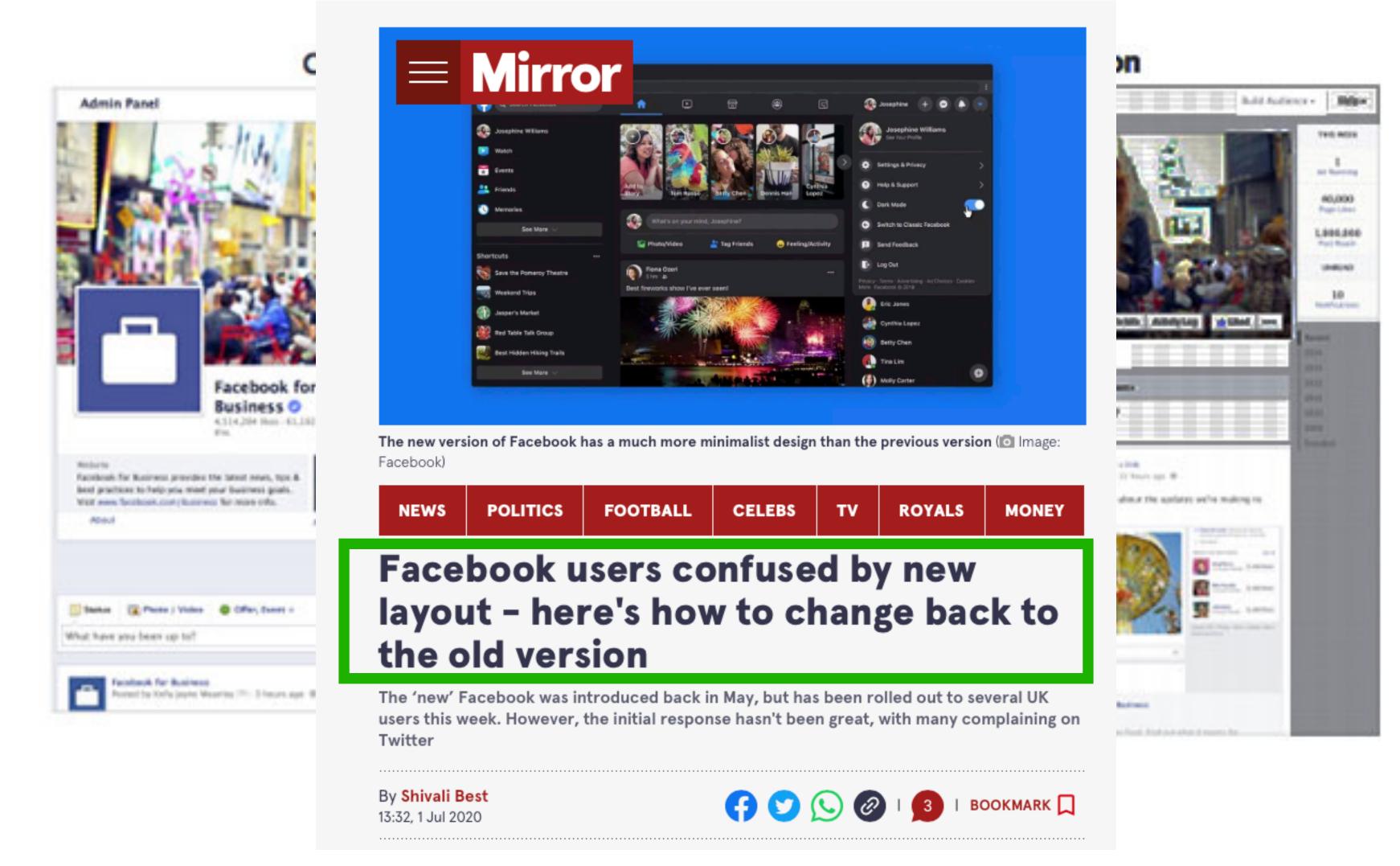
### Example: A/B Testing



## Example: A/B Testing



## Example: A/B Testing



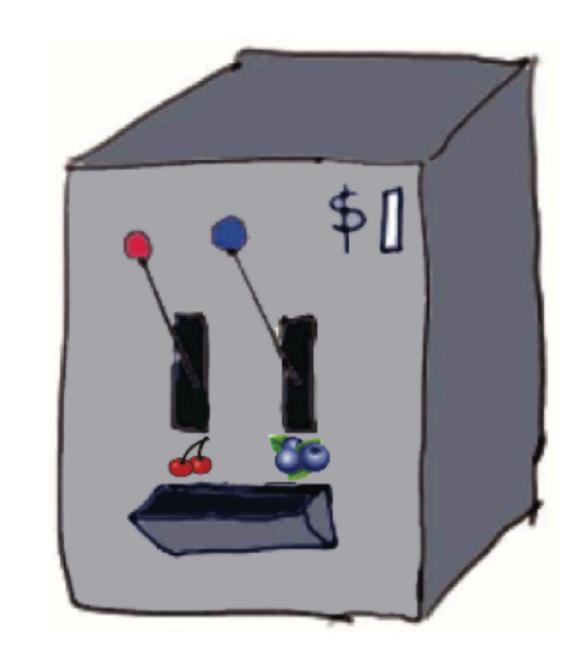
Kimang KHUN

Exploration-Exploitation

#### What can a scientist do?

We need mathematical formulation:

- 1. to design new models (algorithms)
- 2. to quantify the trade-off

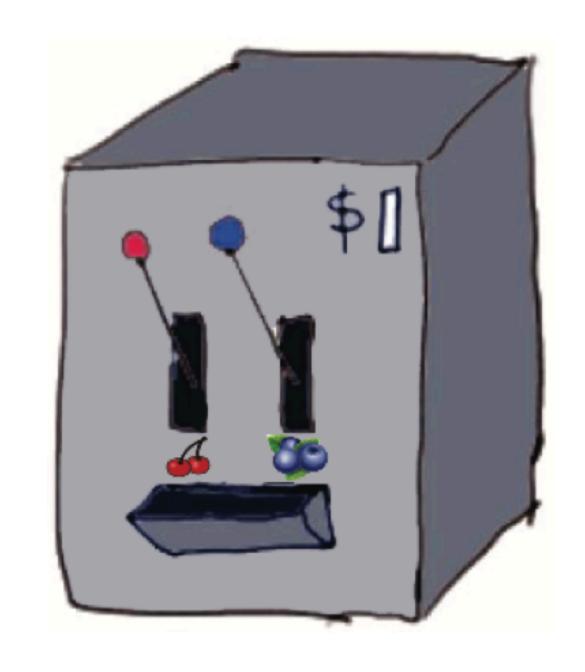


Arm 1 Arm 2  $\mu_1$ 

Pulling Arm 1 gives 1 w.p.  $\mu_1$  or 0 w.p.  $1 - \mu_1$ .

A possible sequence of outcomes: 1,0,0,1,1,... s.t.

$$\mu_1 = \lim_{T \to \infty} \frac{1}{T} \left( \underbrace{1 + 0 + 0 + 1 + 1 + \dots}_{T \text{ terms}} \right)$$



Arm 1 Arm 2  $\mu_1$   $\mu_2$ 

Pulling Arm 2 gives 1 w.p.  $\mu_2$  or 0 w.p.  $1 - \mu_2$ . A possible sequence of outcomes: 0,0,0,0,1,... s.t.

$$\mu_2 = \lim_{T \to \infty} \frac{1}{T} \left( \underbrace{0 + 0 + 0 + 0 + 1 + \dots}_{T \text{ terms}} \right)$$

T terms

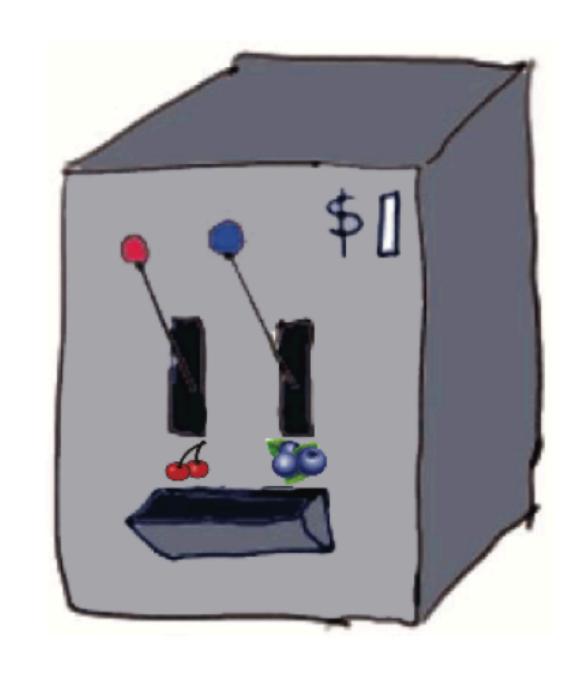


1, 0, 0, 1, 1, 0, 0, 0, 0, 1, ...

Cumulative reward := 1+0+0+1+1+0+0+0+0+1+...

**Question**: Which arm to pull so that the expected cumulative reward is **maximized**?

Arm 1 Arm 2  $\mu_1$   $\mu_2$ 



Arm 1 Arm 2  $\mu_1$   $\mu_2$ 

Question: Which arm to pull so that the expected cumulative reward is maximized?

If  $\mu_1$  and  $\mu_2$  are KNOWN, then

- always pull Arm 1 if  $\mu_1 > \mu_2$
- always pull Arm 2 otherwise.

Challenge:  $\mu_1$  and  $\mu_2$  are UNKNOWN.

This is called "Stochastic bandit".

#### Motivation

#### Maximize clicks

Title	Click probability
"Murder Victim found in an Adult Entertainment Venue"	$\mu_1$
"Headless body found in Topless bar"	$\mu_2$

Choose which title to display. Observe "Click/Not Click".

#### Clinical trials $\mu$



 $\mu_2$ 



Choose treatment for patient.

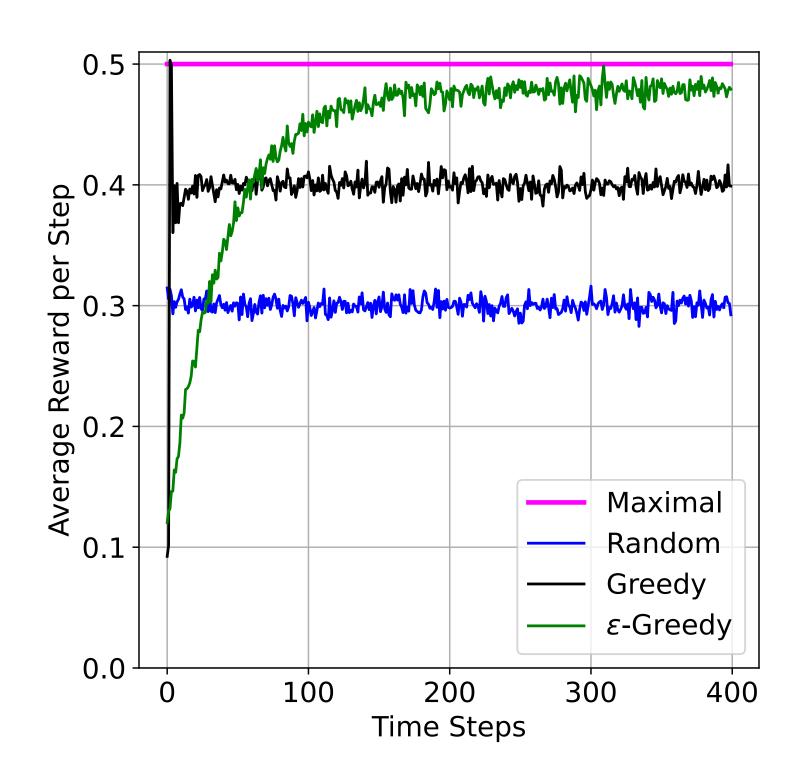
Observe "Heal/Not Heal".

## Algorithm Design

- Random: at each decision time, uniformly randomly pull one arm.
- Greedy: initially try each arm the same number of pulls, then commit on the best arm.
- $\varepsilon$ -Greedy: w.p.  $\varepsilon$ , uniformly randomly pull one arm (Exploration), and w.p.  $1 \varepsilon$ , pull the best arm so far (Exploitation).

## Algorithm Design

- Setup:  $\mu_1 = 0.1$  and  $\mu_2 = 0.5$
- Greedy: try each arm 2 pulls before committing
- $\varepsilon$ -Greedy:  $\varepsilon = 0.1$



Kimang KHUN
Algorithm Design

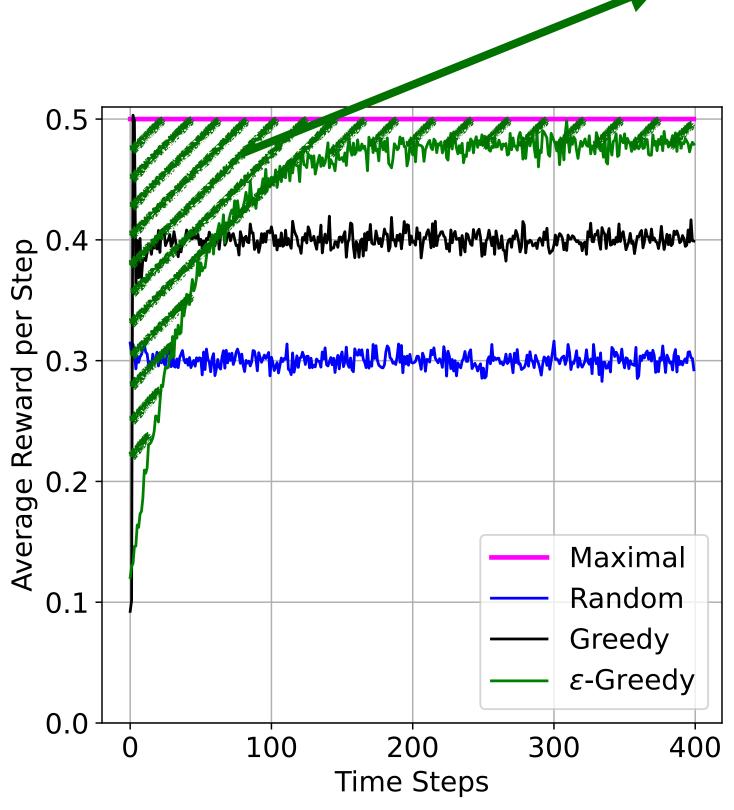
# Algorithm Design

- Setup:  $\mu_1 = 0.1$  and  $\mu_2 = 0.5$ 

- Greedy: try each arm 2 pulls before committing

-  $\varepsilon$ -Greedy:  $\varepsilon = 0.1$ 

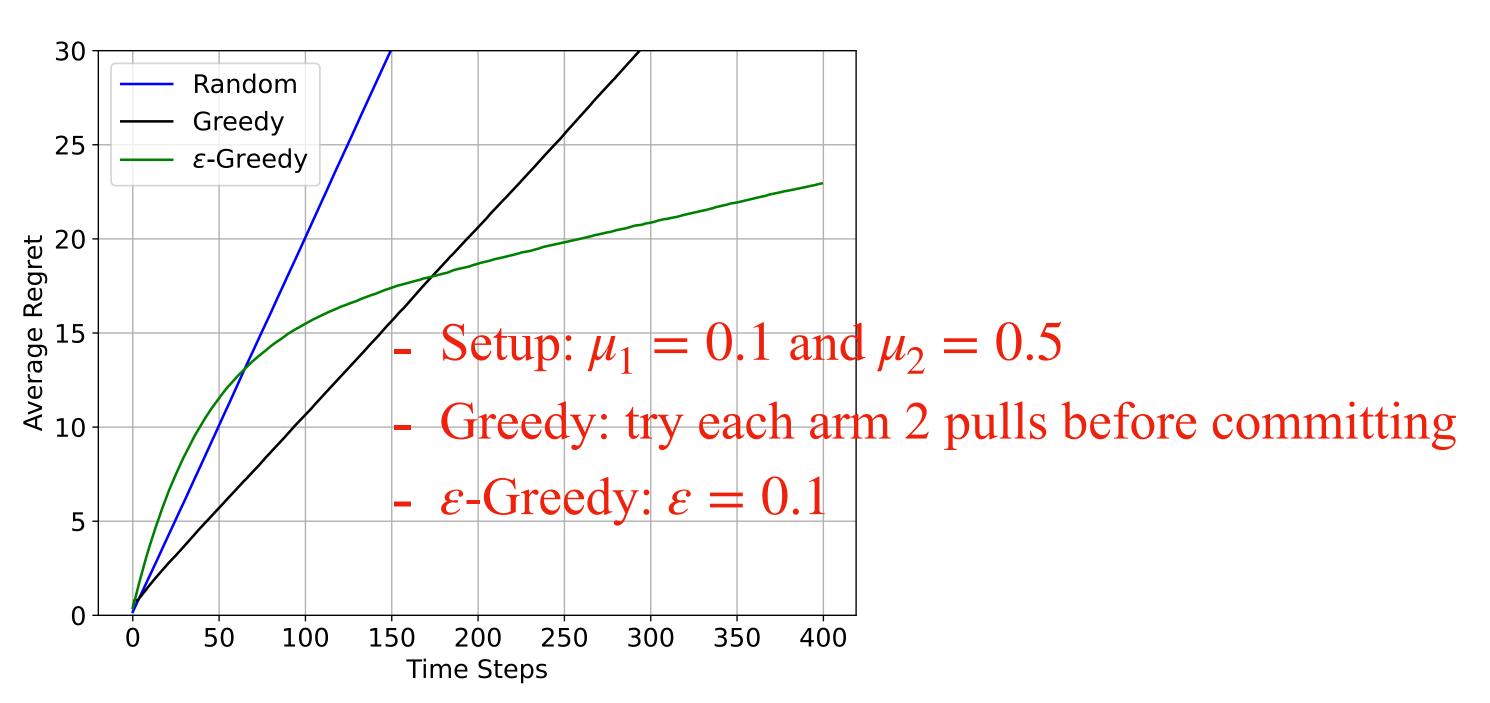
Regret of  $\varepsilon$ -Greedy



### Performance metric: Regret

Regret of  $\mathcal{A} := (\underline{\text{maximal}} \text{ cumulative reward})$  - (cumulative reward of  $\mathcal{A}$ ).

The smaller the regret is, the better  $\mathcal{A}$  performs.



### Performance metric: Regret

Regret of  $\mathcal{A} := (\underline{\text{maximal}} \text{ cumulative reward})$  - (cumulative reward of  $\mathcal{A}$ ).

The smaller the regret is, the better  $\mathcal{A}$  performs.

Let T be total steps.

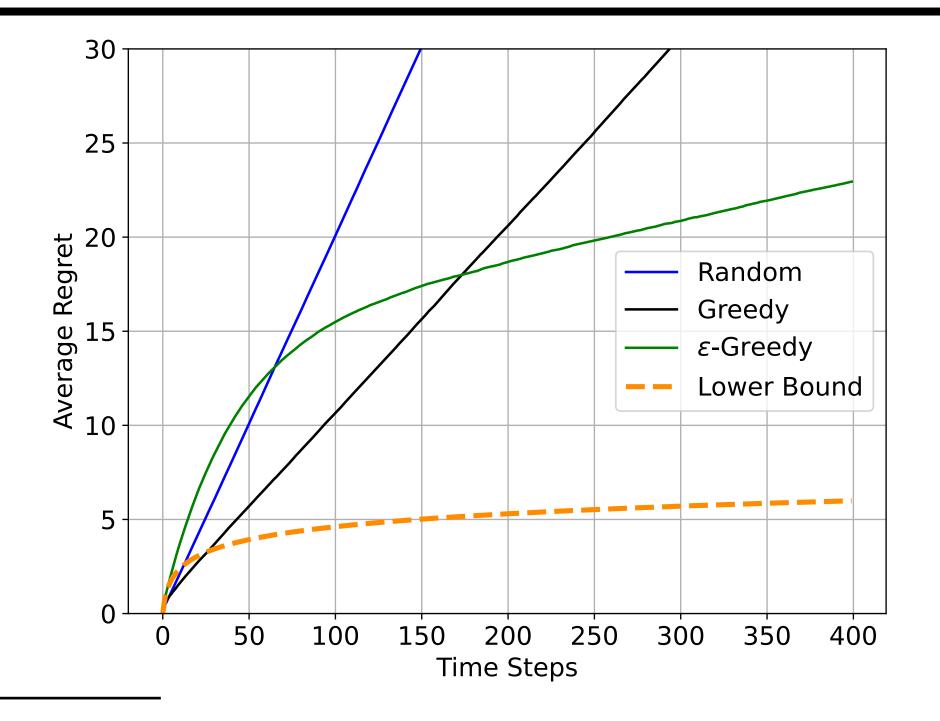
The regret of  $\varepsilon$ -Greedy is O(T) (this is called linear regret).

Can we do better?

### Lower bound on Regret

Theorem 1 (Lai & Robbins, 1985)

There exists a constant c (that depends on  $\mu$ ) s.t. any uniformly efficient algorithm  $\mathcal{A}$  satisfies:  $Regret\ of\ \mathcal{A} \ge c \ln T.$ 

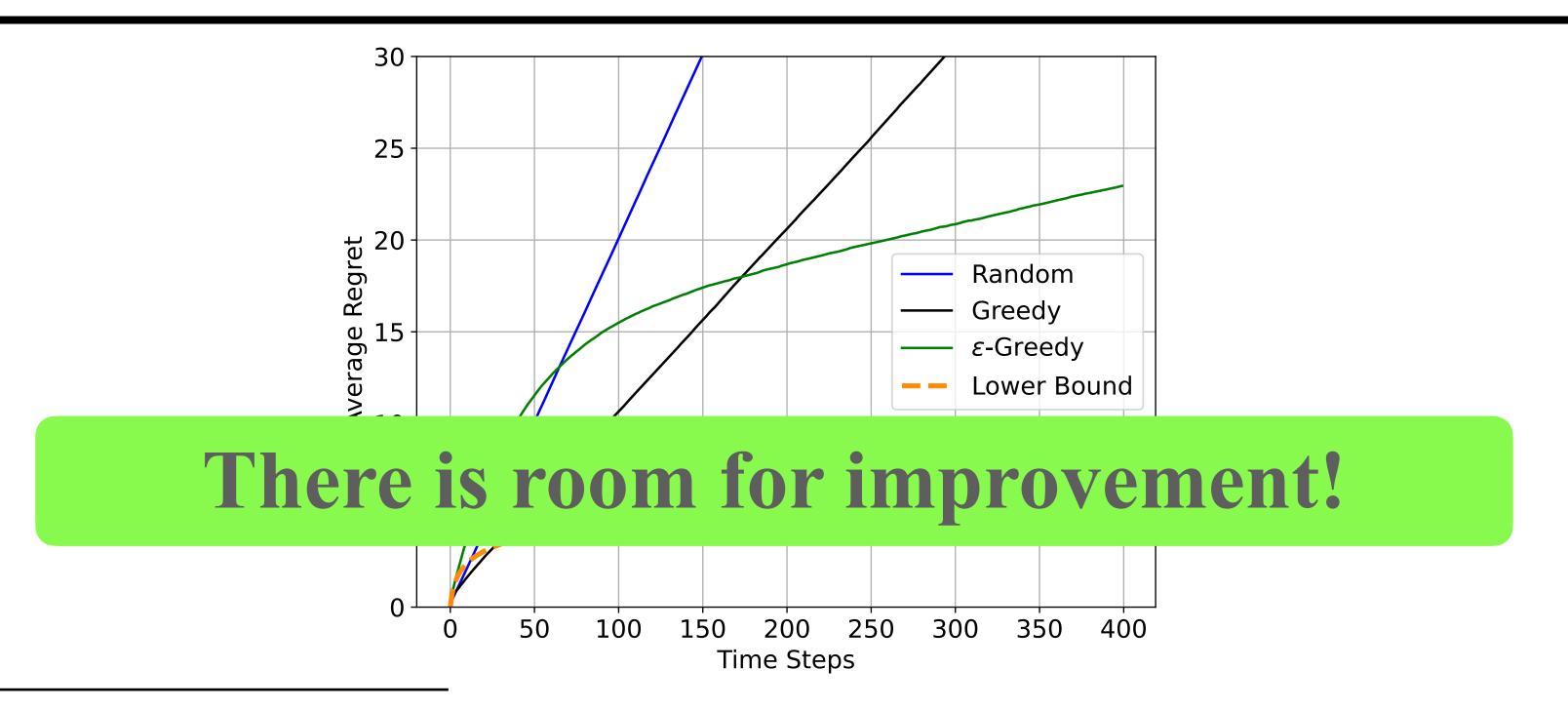


<sup>&</sup>lt;sup>1</sup>Meaning that Regret of  $\mathscr{A}$  is  $o(T^{\alpha})$  for all  $\mu$  and  $\alpha$ .

## Lower bound on Regret

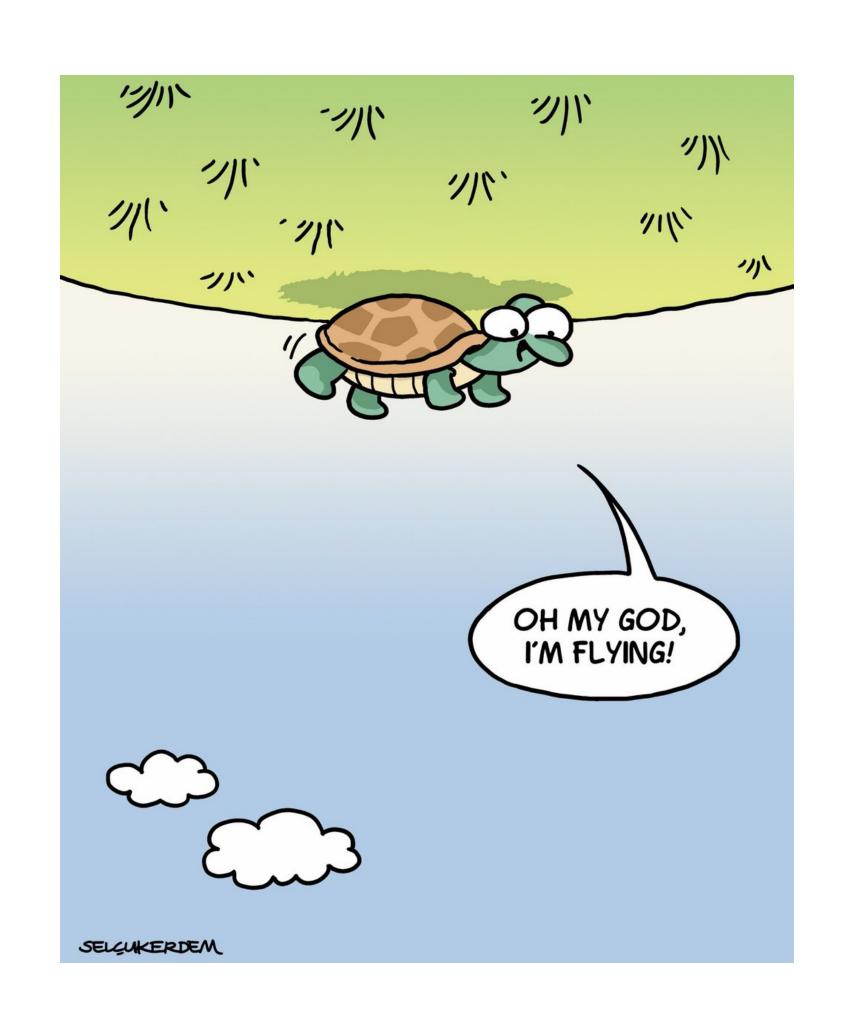
Theorem 1 (Lai & Robbins, 1985)

There exists a constant c (that depends on  $\mu$ ) s.t. any uniformly efficient algorithm  $\mathcal{A}$  satisfies:  $Regret\ of\ \mathcal{A}\ \geq c\ln T.$ 



<sup>&</sup>lt;sup>1</sup>Meaning that Regret of  $\mathscr{A}$  is  $o(T^{\alpha})$  for all  $\mu$  and  $\alpha$ .

## Optimism in Face of Uncertainty (OFU)



When you are uncertain, consider the best possible environment.

If the best possible environment is correct

⇒ No reward lost

Exploitation

If the best possible environment is wrong

⇒ Gather useful info.

Exploration

Consider a coin that gives "Head" w.p.  $\mu$ .

Suppose that you toss the coin N times and observe "Head" n times.

The natural estimator of  $\mu$  is:

$$\hat{\mu} := \frac{n}{N}$$

By Hoeffding's inequality, we have that  $^2$  for x > 0,

$$\mathbb{P}\left\{-\sqrt{\frac{x}{2N}} + \hat{\mu} \le \mu \le \hat{\mu} + \sqrt{\frac{x}{2N}}\right\} \ge 1 - 2e^{-x}.$$

<sup>&</sup>lt;sup>2</sup>under the assumption that all the observations are i.i.d.

Consider a coin that gives "Head" w.p.  $\mu$ .

Suppose that you toss the coin N times and observe "Head" n times.

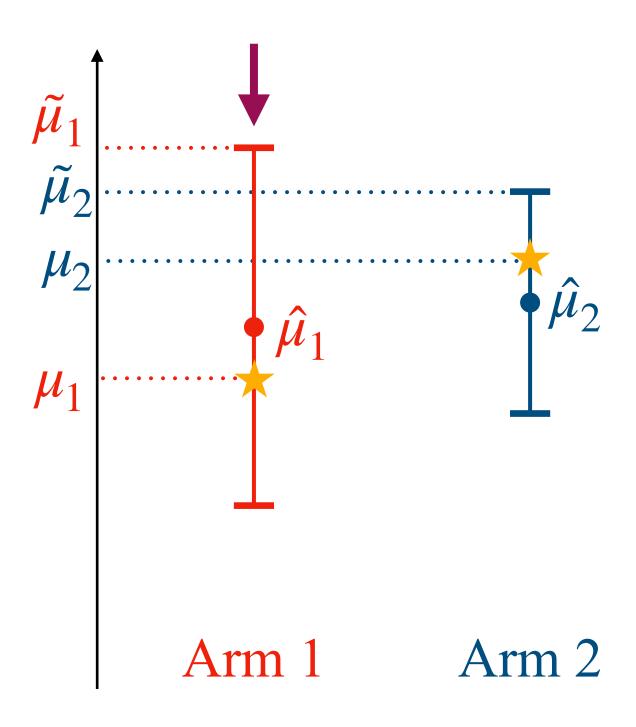
The natural estimator of  $\mu$  is:

$$\hat{\mu} := \frac{n}{N}$$

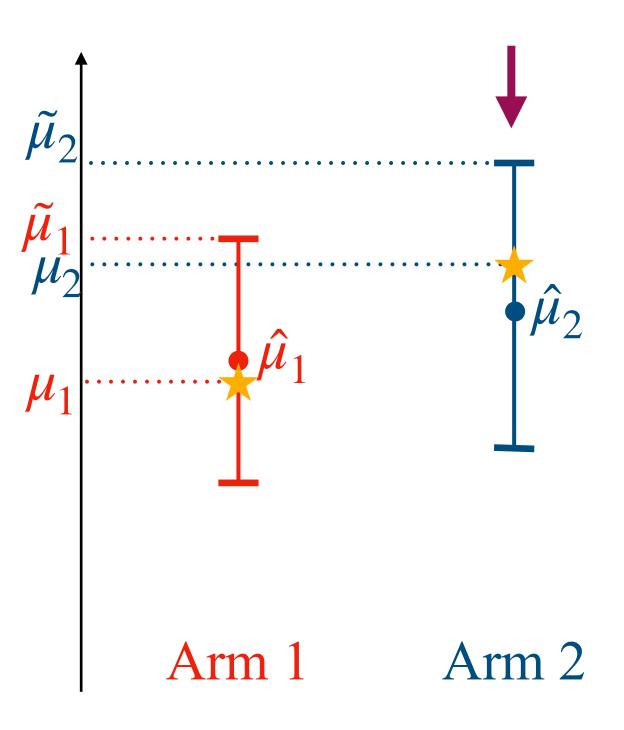
By Hoeffding's inequality, we have that  $^2$  for x > 0,

$$\mathbb{P}\left\{-\sqrt{\frac{x}{2N}} + \hat{\mu} \le \mu \le \hat{\mu} + \sqrt{\frac{x}{2N}}\right\} \ge 1 - 2e^{-x}.$$

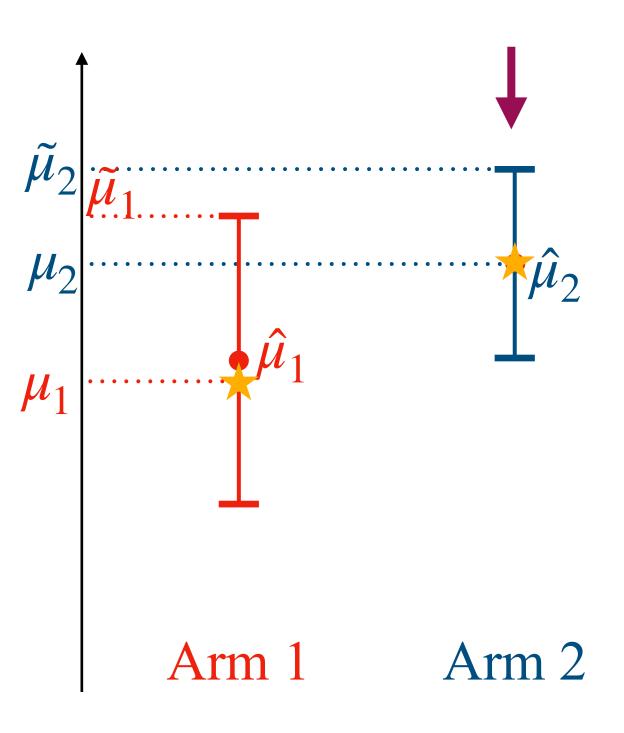
<sup>&</sup>lt;sup>2</sup>under the assumption that all the observations are i.i.d.



At time step t

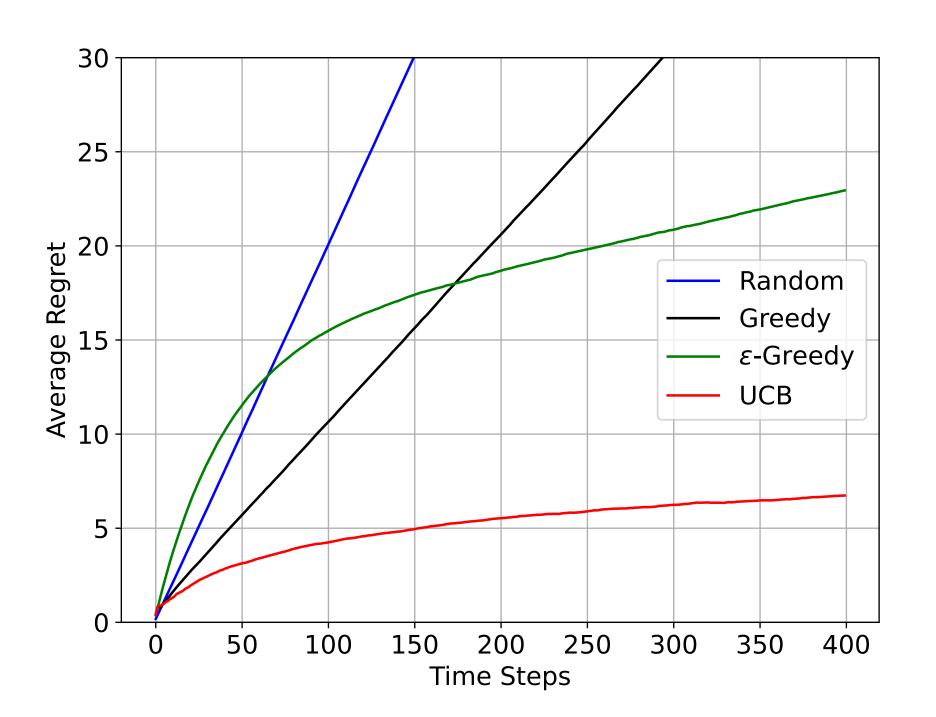


At time step t + 1



At time step t + 2

- Setup:  $\mu_1 = 0.1$  and  $\mu_2 = 0.5$
- Greedy: try each arm 2 pulls before committing
- $\varepsilon$ -Greedy:  $\varepsilon = 0.1$



Theorem 2 (Auer et al., 2002):

Regret of UCB  $\leq c' \ln T$ .

We say that UCB is asymptotically optimal.

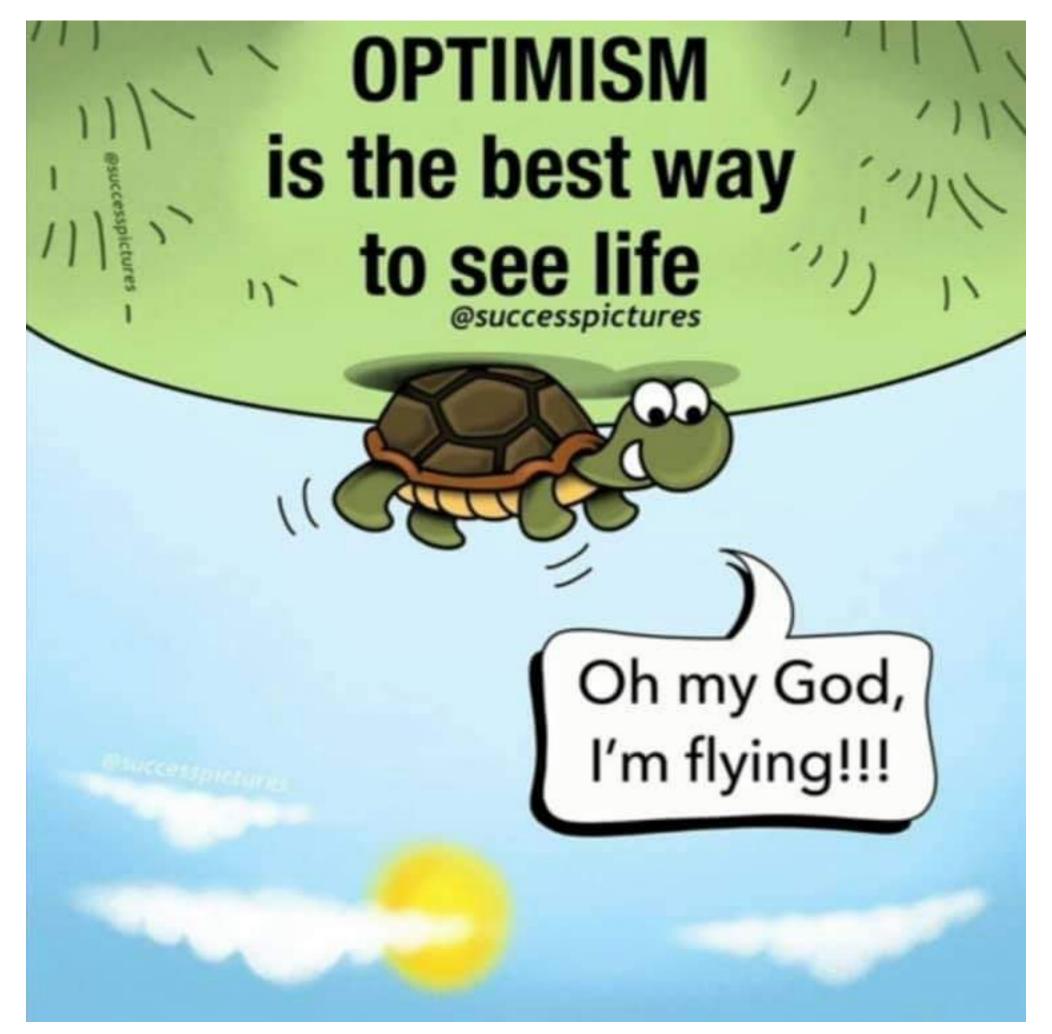
Kimang KHUN
Algorithm Design

#### Conclusion

- Exploration vs. Exploitation (EE) dilemma happens in the world of decision-making under uncertainty.
- Multi-armed bandit (MAB) problem is a mathematical formulation allowing us to consider the EE trade-off and design new algorithms.
- We use Regret to measure the algo.'s performance.
- No algorithm has a regret smaller than  $O(\ln T)$  uniformly over all MAB problems.
- UCB algorithm from OFU approach has a regret bounded by  $O(\ln T)$  (it is asymptotically optimal).

For more on bandit, check out this book





source: https://twitter.com/parveenkaswan/status/1364791588442890240?lang=zh-Hant

https://kimang18.github.io or khun.kimang@misti.gov.kh

Questions?

Kimang KHUN

Conclusion